

# Finding Optimal Size TDMA Schedules using Integer Programming

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## Abstract

The problem of finding minimal size TDMA schedules is formally described as an Integer Program (IP). A brief user manual explains how the attached implementation can be used to find minimal size TDMA for any given WSN and routing table, fulfilling the validity criteria.

## 1 TDMA scheduling

The TDMA problem considers a given routing for a (weighted) Wireless Sensor Network (WSN) with a  $n \times n + q$  connectivity matrix  $c$  and a  $n \times n + q$  routing matrix  $r$ , with  $n$  being the amount of sending sensors, and  $q$  the amount of sinks. The problem is formulated as the maximization of  $f(c, r) = \sum_i E_i, i \in \{1, \dots, m\}, E_i \in \{0, 1\}$ , the amount of empty slots, given the upper bound

$$m = \sum_{k=1}^n k = \frac{n(n+1)}{2}$$

for the case that each sensor has one packet waiting for delivery.

The formulation can be found in Table 1. It is restricted to single parent routing (assuming r-r conflict free schedules), and expects the given graph  $r$  to be a strict order with the sinks being modeled in position sensor  $n + 1, \dots, n + q$ .  $E_i$  is the amount of empty slots in a TDMA of length  $m$ . The formulation maximizes  $E$ , meaning that it seeks to find a slot table with a maximum of empty numbers which is equal to finding the shortest TDMA. The result is the  $m \times n$  matrix  $S$  (optimal size TDMA).

maximize	$\sum_i E_i$	amount of empty slots
subject to	$\sum_j S_{ij} - Y_{1i} \leq 0$ $E_i - (1 - Y_{1i}) \leq 0$	c0. empty slot definition ( $\sum_j S_{ij} > 0 \Rightarrow E_i = 0$ )
	$B_{1j} = 1$ $B_{ij} \geq 0$	init: buffers start with 1 waiting packet buffers are non-negative
	$S_{ij} - amountSensors * Y_{2ijl} \leq 0$ $c_{lj} S_{il} - amountSensors * (1 - Y_{2ijl}) \leq 0$	c1. t-r, r-t-r conflicts ( $S_{ij} = 1 \Rightarrow c_{lj} S_{il} = 0$ )
	$\sum_j S_{ij} * C_{jl} \leq 1$	c2. t-t, t-r-t conflicts
	$-\sum_j S_{ij} - amountSensors * Y_{3i} \leq -1$ $\sum_j B_{ij} - amountSensors * (1 - Y_{3i}) \leq 0$	all buffers are empty after last delivery
	$\sum_j R_{ij} = 1$	single parent routing
	$S_{ij} - amountSensors * Y_{4ij} \leq 0$ $-B_{ij} - amountSensors * (1 - Y_{4ij}) \leq -1$	empty buffer $\rightarrow$ no submission
	$B_{ij} = B_{(i-1)j} - S_{(i-1)j} + \sum_l X_{ijl}$ $R_{lj} + S_{(i-1)l} - Y_{1ijl} \leq 1$ $X_{ijl} + (1 - Y_{1ijl}) \geq 1$	buffer depends on former transitions $R_{il} + S_{(i-1)l} = 2 \Leftrightarrow X_{ijl} = 1$ slack variable X models arrivals
	$X_{ijl} - Y_{2ijl} \leq 0$ $R_{lj} + S_{(i-1)l} + 2 * (1 - Y_{2ijl}) \geq 2$	
	$X_{ijl} + Y_{3ijl} \geq 1$ $R_{lj} + S_{(i-1)l} - (1 - Y_{3ijl}) \leq 1$	

Table 1: Parameters are written with small letters, variables (to be optimized) with capital letters. The formulation could possibly be improved with regards to efficiency.

## 2 Program Description

In order for the program to work you have to have Java JRE 1.5 and the gurobi solver in version 5 or higher installed <sup>1</sup>. Once you have registered your installation and set your gurobi environment variables appropriately, you can start the `tdmaSolver.jar` by:

- `java -jar tdmaSolver.jar.`

The program expects a properties file by name `'wsnSetup'` in the same folder, which defines three parameters:

**wsn** the path to the network to be solved (mandatory)

**routing** the path to a routing table for the network (optional)

**buffer** the initial buffer setting for the network (optional)

The file content could look as in:

<sup>1</sup>Take a look at [www.gurobi.com](http://www.gurobi.com) for an howto on installing it (you can attain a academic license for scientific purposes)

```

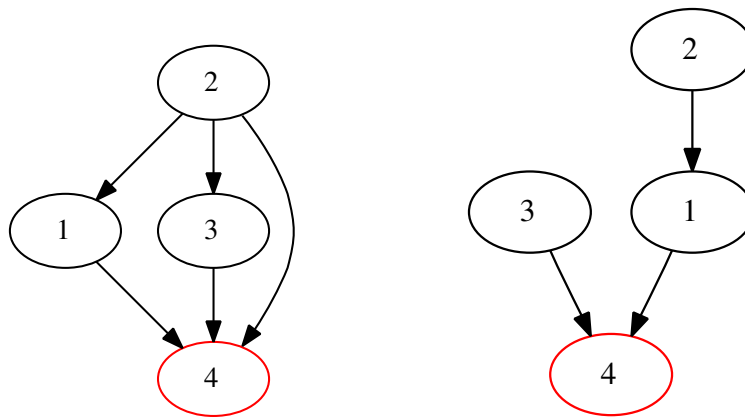
wsn=myWSN.txt
routing=myRouting.txt
buffer=myBuffer.txt

```

Solely defining the network potentially leads to better results (and substantially longer running times), because of less restrictions (no extra routing constraints). The upper bound  $m$  takes the amount of initial packets in the buffer into consideration.

### 3 Example

Given are the network in Figure 1a and the valid routing 1b.



(a) A trivial example WSN with 4 nodes. (b) A routing for the trivial example WSN in 1a.

The equivalent in-file description is formulated as in:

```

WSN:
00000000100000000000
00100000001000000010
00010001000000000100
00100000000010000000
0000000000000000010000
1000000000000000000000
0100000000000000000000
000000000010000000001
00000000000000000010000
00100000000000000000000
00010000000000000100000
010000001000010000000
01000000001000010000

```

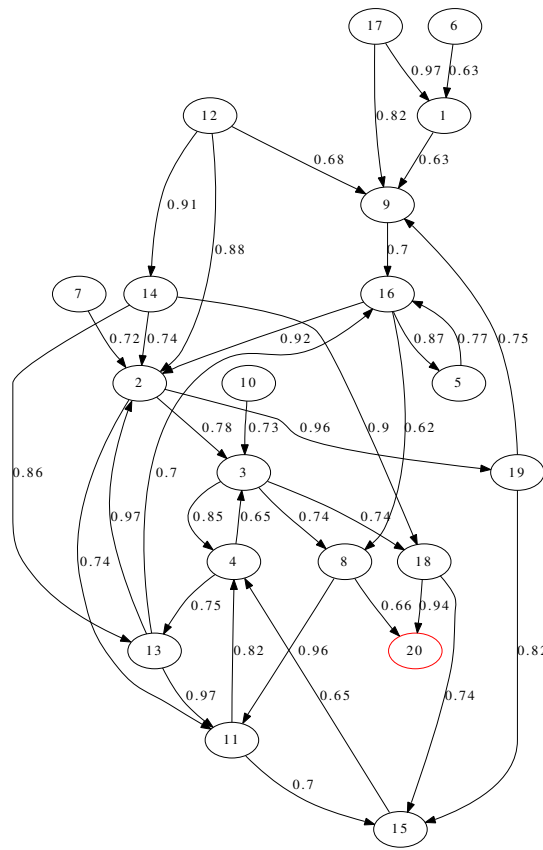


Figure 1: The test network to be optimized, consisting of 20 nodes.

```

01000000000010000100
00010000000000000000
01001001000000000000
10000000100000000000
00000000000000100001
00000000100000100000

```

1s stand for positive connectivity, 0s for no connectivity in the WSN file. In the routing file, 0s stand for no route, 1s for given route. Currently, based on the former problem formulation and this syntax, only non-weighted graphs are supported for optimization. The program will print the length of the optimal TDMA via system out, followed by one<sup>2</sup> optimal size schedule. Important note:

- **The algorithm gives only appropriate solutions for graphs that connect each sensor to a sink!**

The output for the example looks as follows:

```

Optimize a model with 553501 rows, 506880 columns and 1249270 nonzeros
Presolve removed 410499 rows and 458901 columns (presolve time = 6s) ...
Presolve removed 410499 rows and 458901 columns (presolve time = 10s) ...
Presolve removed 461604 rows and 458903 columns
Presolve time: 11.61s
Presolved: 91897 rows, 47977 columns, 324651 nonzeros
Variable types: 0 continuous, 47977 integer (45621 binary)
Found heuristic solution: objective -0.0000000

```

Root simplex log...

Iteration	Objective	Primal Inf.	Dual Inf.	Time
6448	1.2025872e+02	1.353778e+04	0.000000e+00	5s
9200	1.2025797e+02	1.434810e+04	0.000000e+00	10s
11264	1.2025670e+02	3.295059e+05	0.000000e+00	15s
12812	1.2024009e+02	1.500940e+07	0.000000e+00	20s
14360	1.1783776e+02	1.785668e+06	0.000000e+00	25s
15736	1.1633980e+02	1.913384e+05	0.000000e+00	30s
17112	1.1570554e+02	3.282639e+06	0.000000e+00	36s
18488	1.1508934e+02	4.145550e+05	0.000000e+00	40s
Warning: Markowitz tolerance tightened to 0.03125				
19864	1.1433710e+02	7.922502e+05	0.000000e+00	45s
21412	1.1411626e+02	1.050797e+05	0.000000e+00	50s
22788	1.1380968e+02	8.141415e+06	0.000000e+00	55s
23992	1.1367013e+02	1.728571e+05	0.000000e+00	60s
25368	1.1348785e+02	6.207536e+04	0.000000e+00	65s
26744	1.1324074e+02	5.421782e+04	0.000000e+00	70s

<sup>2</sup>There could be multiple, but the program stops after finding one.

28120	1.1317063e+02	1.029490e+04	0.000000e+00	75s
29324	1.1311170e+02	9.164295e+03	0.000000e+00	80s
30872	1.1298886e+02	6.719229e+02	0.000000e+00	85s
32420	1.1293671e+02	4.795785e+04	0.000000e+00	91s
33624	1.1293410e+02	3.400464e+03	0.000000e+00	95s
34353	1.1293510e+02	0.000000e+00	0.000000e+00	98s

Root relaxation: objective 1.129351e+02, 34353 iterations, 98.18 seconds  
Total elapsed time = 115.54s

Nodes	Current Node			Objective Bounds		Gap	Work		
	Expl	Unexpl	Obj Depth IntInf	Incumbent	BestBd		It/Node	Time	
0	0	112.93510	0 3339	-0.00000	112.93510	-	-	118s	
0	0	99.38295	0 3190	-0.00000	99.38295	-	-	271s	
0	0	99.22089	0 3338	-0.00000	99.22089	-	-	288s	
0	0	99.16583	0 3339	-0.00000	99.16583	-	-	305s	
0	0	99.13332	0 3372	-0.00000	99.13332	-	-	315s	
0	0	99.11349	0 3415	-0.00000	99.11349	-	-	324s	
0	0	99.10528	0 3342	-0.00000	99.10528	-	-	332s	
0	0	99.10147	0 3314	-0.00000	99.10147	-	-	563s	
0	0	99.09358	0 3382	-0.00000	99.09358	-	-	568s	
0	0	99.08600	0 3373	-0.00000	99.08600	-	-	574s	
0	0	99.08265	0 3382	-0.00000	99.08265	-	-	580s	
0	0	99.07641	0 3183	-0.00000	99.07641	-	-	592s	
0	0	99.06997	0 3322	-0.00000	99.06997	-	-	605s	
0	0	99.06845	0 3247	-0.00000	99.06845	-	-	619s	
0	0	99.06437	0 3324	-0.00000	99.06437	-	-	630s	
0	0	99.05293	0 3315	-0.00000	99.05293	-	-	644s	
0	0	99.05145	0 3469	-0.00000	99.05145	-	-	653s	
0	0	99.04969	0 3270	-0.00000	99.04969	-	-	663s	
0	0	99.04928	0 3314	-0.00000	99.04928	-	-	675s	
0	0	99.04839	0 3294	-0.00000	99.04839	-	-	688s	
0	0	99.04681	0 3338	-0.00000	99.04681	-	-	699s	
0	0	99.04662	0 3290	-0.00000	99.04662	-	-	711s	
0	0	99.04589	0 3357	-0.00000	99.04589	-	-	721s	
0	0	99.00354	0 3361	-0.00000	99.00354	-	-	732s	
0	0	99.00354	0 3360	-0.00000	99.00354	-	-	734s	
H	0			94.0000000	99.00354	5.32%	-	782s	
	0	2	99.00354	0 3360	94.00000	99.00354	5.32%	-	791s
	1	3	98.92936	1 3341	94.00000	98.92936	5.24%	2901	795s
	3	3	cutoff	3	94.00000	98.82752	5.14%	7364	878s
	4	4	98.73492	3 3238	94.00000	98.73492	5.04%	5848	883s
	5	2	cutoff	4	94.00000	98.73126	5.03%	6550	925s
	6	3	97.93249	4 3145	94.00000	97.93249	4.18%	7461	988s
	8	1	cutoff	5	94.00000	97.89818	4.15%	6423	1016s

9	2	97.09808	5 3215	94.00000	97.09808	3.30%	6102	1031s
12	2	cutoff	7	94.00000	97.09642	3.29%	6727	1186s
13	3	96.82057	6 3132	94.00000	96.97763	3.17%	6741	1214s
14	4	96.23050	7 2939	94.00000	96.82057	3.00%	6452	1223s
15	3	cutoff	8	94.00000	96.82057	3.00%	7037	1299s
16	4	95.33293	8 2944	94.00000	96.82057	3.00%	6696	1304s
17	2	95.33293	9 2944	94.00000	96.82057	3.00%	6302	1345s
21	3	95.06702	9 2950	94.00000	96.82057	3.00%	5657	1366s

Cutting planes:

Gomory: 4  
 Cover: 980  
 Implied bound: 1365  
 Clique: 1301  
 MIR: 5  
 GUB cover: 166  
 Zero half: 27

Explored 26 nodes (258052 simplex iterations) in 1367.17 seconds  
 Thread count was 2 (of 2 available processors)

Optimal solution found (tolerance 1.00e-04)  
 Best objective 9.400000000000e+01, best bound 9.400000000000e+01, gap 0.0%  
 Routing:

A result of minimal size 31 could be found in 1367406 ms.

Routing:

00000000100000000000  
 00000000000000000010  
 00010000000000000000  
 00000000000010000000  
 0000000000000010000  
 10000000000000000000  
 01000000000000000000  
 00000000000000000001  
 000000000000000010000  
 00100000000000000000  
 00010000000000000000  
 00000000000001000000  
 01000000000000000000  
 00000000000000000100  
 00010000000000000000  
 00000001000000000000  
 10000000000000000000  
 00000000000000000001

000000010000000000

Schedule and buffer:

0100000111100000000	1111111111111111111
0000000000001111001	1022111000011112112
1001000000000001010	1123111110010001121
0001010110000000000	0122111220011000111
0010000110000000100	1121101110012001111
1010001000000001000	2112101000012002011
1000100100000000000	1203100110012001011
0000000000001001010	0203000020012002011
0100000110000000000	0303000120011001001
01000000000000001000	0203000010011002002
0101000110000000000	0103000110011001003
0000000000001001000	000200000012002004
0001000000010001001	0102000100011001004
0000000110000000000	0101000210002100003
0000000100001000001	0101000100002101003
0000000000001001001	0201000010001101002
0101000110000100000	0301000120000100001
01000000000000001010	0200000010001001012
0100000110000000000	0100000110001000003
00000000000000001001	0000000000001001004
0000000110001000000	0000000110001000003
00000000000000001001	0100000000000001003
0100000110000000000	0100000110000000002
00000000000000001001	0000000000000001003
0000000110000000000	0000000110000000002
00000000000000001001	0000000000000001002
0000000110000000000	0000000110000000001
00000000000000001001	0000000000000001001
0000000110000000000	0000000110000000000
00000000000000001000	0000000000000001000
0000000100000000000	0000000100000000000